Complexity Analysis in Cyclic Tag System Emulated by Rule 110

AUTOMATA2013

Sep. 17 – 19, 2013

Giessen, Germany

Shigeru Ninagawa ^{1,3} and Genaro J. Martínez ^{2,3} ¹ Kanazawa Institute of Technology, Japan ² Instituto Politécnico Nacional, Mexico ³ University of the West of England, United Kingdom

Table of contents

- 1. Background
- 2. Lempel-Ziv complexity
- 3. Cyclic tag system emulated by rule 110
- 4. Results
- 5. Conclusion

Basic idea

- CAs have no memory except for cell
 → all the information necessary to
 perform computation is in its configuration
- Complexity of conf. is reflective of the complexity of information necessary for computation

Related works

- Compression-based classification of ECA by DEFLATE (LZ77 + Huffman coding) (Zenil, 2010)
- Parity problem solving process by ECA rule 60 in array size of 2ⁿ (Ninagawa, 2012) ← stepwise decrease by period halving (Lempel-Ziv complexity, LZ78)

Motivation

- ECA rule 110 is supporting universal computation by emulating cyclic tag system
- How does complexity vary during the emulation process by rule 110?
- We employ Lempel-Ziv complexity as a measure of complexity

Lempel-Ziv complexity (Ziv, Lempel, 1978)(1/2)

 $s_1s_2 \cdots s_ks_{k+1} \cdots$: given string, $s_i \in$ alphabet $s_1 \cdots s_k$ has already been divided into phrases $w_1 \cdots w_m (= s_1s_2 \cdots s_k), m \leq k, w_0 = \epsilon$ (empty string) search the longest substring $s_{k+1} \cdots s_{k+n} = w_j \ (0 \leq j \leq m)$ and set $w_{m+1} = w_j s_{k+n+1}$



Lempel-Ziv complexity (2/2)

For example: $010010101 \cdots$ is given $w_0 = \epsilon$, $w_1 = 0 = w_00$, $w_2 = 1 = w_01$, $w_3 = 00 = w_10$, $w_4 = 10 = w_20$, $w_5 = 101 = w_41$, ...

The number of divided phrases: Lempel-Ziv complexity

Cyclic tag system(Cook, 2004)

- Σ={0,1}, tape is read from the front and appended according to appendant table
- Example appendant table (1, 101)
 - 1
 appendant

 11
 1

 1101
 101

 1011
 1

 011
 101 skipped in reading '0'

 111
 1
- CTS can emulate tag system and rule 110 can emulate CTS

Basic mechanism (tape data '1')



Space-time pattern of ossifier



Basic mechanism (tape data '1')





Basic mechanism (tape data '1')





Collision between ossifier and moving data creates tape data

Basic mechanism (tape data '0')



O · · · appendant X skipped



Rejector is erasing table data

IC emulating CTS (N=65,900)

http:ucomp.uwe.ac.uk/genaro/rule110/ctsRule110.html

Left edge (0)

Ossifier * 6 (7230 - 50100)

Tape data'1' (52600 - 52750)

Table data '1' Leader2 Leader1 Table data '101' (53100 - 53600)(53600 - 53900)(52750-53100)(53900 - 55150)Leader3 Table data '1' Leader4 Table data '101' (55200 - 55500)(55500 - 55950)(55950 - 56200)(56200 - 57550)Leader5 Table data '1' Leader6 (57550 - 57900)Right edge (57900 - 58350)(58350 - 58700)(65899)



Evolution of LZ complexity



Array size: 65,900

Enlarged view of LZ complexity in cts emulation



Moving average of LZ complexity in CTS emulation (period: 100)



Evolution of LZ complexity in each section

whole array is divided into 20 sections (3,295 cells each) Sec. 0: leftmost, sec.19: rightmost



Moving average of LZ complexity in the three parts (array size: 1,100) of sec.14

Moving average of LZ complexity

Table data '1'

Moving data '1'

Conclusion

- In the emulation process of CTS by rule 110, stepwise decrease of LZ complexity is observed
- When table data are transformed into moving data by acceptor or erased by rejector, LZc decreases quickly
- These results might generally hold for decision problem solving process.

Thank you for listening!

X=43000..52999

